Numerical Analysis - Part II

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Lecture 17

Iterative methods for linear algebraic systems

Solving linear systems with iterative methods

The general *iterative* method for solving $Ax = b$ is a rule $\mathbf{x}^{k+1} = f_k(\mathbf{x}^0, \mathbf{x}^1, \ldots, \mathbf{x}^k)$. We will consider the simplest ones: *linear, one-step, stationary* iterative schemes:

$$\mathbf{x}^{k+1} = H\mathbf{x}^k + \mathbf{v}, \qquad \mathbf{x}^0, \mathbf{v} \in \mathbb{R}^n. \tag{1}$$

Here one chooses H and v so that x^*, a solution of $Ax = \mathbf{b}$, satisfies $\mathbf{x}^* = H\mathbf{x}^* + \mathbf{v}$, i.e. it is the fixed point of the iteration (1) (if the scheme converges). Standard terminology:

the *iteration matrix* H, the *error* $\mathbf{e}^k := \mathbf{x}^* - \mathbf{x}^k$, the *residual* $\mathbf{r}^k := A\mathbf{e}^k$

Solving linear systems – Iterative refinement

For a given class of matrices A (e.g. positive definite matrices, or even a single particular matrix), we are interested in *convergent* methods, i.e. the methods such that $\boldsymbol{x}^k \to \boldsymbol{x}^* = A^{-1}\boldsymbol{b}$ for every starting value $\boldsymbol{x}^0$. Subtracting $\boldsymbol{x}^* = H\boldsymbol{x}^* + \boldsymbol{v}$ from (1) we obtain

$$\boldsymbol{e}^{k+1} = H\boldsymbol{e}^k = \cdots = H^{k+1}\boldsymbol{e}^0, \tag{2}$$

i.e., a method is convergent if $\boldsymbol{e}^k = H^k\boldsymbol{e}^0 \to 0$ for any $\boldsymbol{e}^0 \in \mathbb{R}^n$.

(Iterative refinement). This is the scheme

$$\boldsymbol{x}^{k+1} = \boldsymbol{x}^k - S(A\boldsymbol{x}^k - \boldsymbol{b}).$$

If $S = A^{-1}$, then $\boldsymbol{x}^{k+1} = A^{-1}\boldsymbol{b} = \boldsymbol{x}^*$, so it is suggestive to choose S as an approximation to A^{-1}. The iteration matrix for this scheme is $H_S = I - SA$.

Solving linear systems – Splitting

(Splitting). This is the scheme

$$(A - B)\mathbf{x}^{k+1} = -B\mathbf{x}^k + \mathbf{b},$$

with the iteration matrix $H = -(A - B)^{-1}B$. Any splitting can be viewed as an iterative refinement (and vice versa) because

$$(A - B)\mathbf{x}^{k+1} = -B\mathbf{x}^k + \mathbf{b} \quad \Leftrightarrow \quad (A - B)\mathbf{x}^{k+1} = (A - B)\mathbf{x}^k - (A\mathbf{x}^k - \mathbf{b})$$

$$\Leftrightarrow \quad \mathbf{x}^{k+1} = \mathbf{x}^k - (A - B)^{-1}(A\mathbf{x}^k - \mathbf{b}),$$

so we should seek a splitting such that $S = (A - B)^{-1}$ approximates A^{-1}.

Solving linear systems – Convergence

Theorem 1

Let $H \in \mathbb{R}^{n \times n}$. Then $\lim\limits_{k \to \infty} H^k \mathbf{z} = 0$ for any $\mathbf{z} \in \mathbb{R}^n$ if and only if $\rho(H) < 1$.

Proof. 1) Let λ be an eigenvalue of (the real) H, real or complex, such that $|\lambda| = \rho(H) \geq 1$, and let $\mathbf{w}$ be a corresponding eigenvector, i.e., $H\mathbf{w} = \lambda\mathbf{w}$. Then $H^k\mathbf{w} = \lambda^k\mathbf{w}$, and

$$\|H^k\mathbf{w}\|_\infty = |\lambda|^k\|\mathbf{w}\|_\infty \geq \|\mathbf{w}\|_\infty =: \gamma > 0. \tag{3}$$

If $\mathbf{w}$ is real, we choose $\mathbf{z} = w$, hence $\|H^k\mathbf{z}\|_\infty \geq \gamma$, and this cannot tend to zero.

If $\mathbf{w}$ is complex, then $\mathbf{w} = \mathbf{u} + \mathrm{i}\mathbf{v}$ with some real vectors $\mathbf{u}, \mathbf{v}$. But then at least one of the sequences $(H^k\mathbf{u}), (H^k\mathbf{v})$ does not tend to zero. For if both do, then also $H^k\mathbf{w} = H^k\mathbf{u} + \mathrm{i}H^k\mathbf{v} \to 0$, and this contradicts (3).
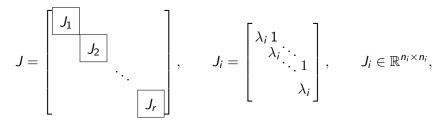
Proof. Cont. 2) Now, let $\rho(H) < 1$, and assume for simplicity that H possesses n linearly independent eigenvectors $(\boldsymbol{w}_j)$ such that $H\boldsymbol{w}_j = \lambda_j \boldsymbol{w}_j$. Linear independence means that every $\boldsymbol{z} \in \mathbb{R}^n$ can be expressed as a linear combination of the eigenvectors, i.e., there exist $(c_j) \in \mathbb{C}$ such that $\boldsymbol{z} = \sum_{j=1}^{n} c_j \boldsymbol{w}_j$. Thus,

$$H^k \boldsymbol{z} = \sum_{j=1}^{n} c_j \lambda_j^k \boldsymbol{w}_j \,,$$

and since $|\lambda_j| \leq \rho(H) < 1$ we have $\lim_{k\to\infty} H^k \boldsymbol{z} = 0$, as required. $\square$

Solving linear systems – Convergence

Remark 2 (Non-examinable)

The complete proof of case (2) of Theorem 1 exploits the so-called Jordan normal form of the matrix H, namely $H = SJS^{-1}$, where J is a block diagonal matrix consisting of the Jordan blocks,

$$
J = \begin{bmatrix} J_1 & & & \\ & J_2 & & \\ & & \ddots & \\ & & & J_r \end{bmatrix}, \qquad J_i = \begin{bmatrix} \lambda_i & 1 & & \\ & \lambda_i & \ddots & \\ & & \ddots & 1 \\ & & & \lambda_i \end{bmatrix}, \qquad J_i \in \mathbb{R}^{n_i \times n_i},
$$

To prove that $J_i^k \to 0$ if $|\lambda_i| < 1$ one should split $J_i = \lambda_i I + P$, notice that $P^m = 0$ for $m \geq n_i$, and evaluate the terms of the expansion $(\lambda_i I + P)^k = \sum_{m=0}^{n_i-1} \binom{k}{m} \lambda_i^{k-m} P^m$.

Solving linear systems – Convergence

Applying Theorem 1 to the error estimate (2), we arrive at the following statement.

Theorem 3
Let $\mathbf{x}^*$, a solution of $A\mathbf{x} = \mathbf{b}$, satisfy $\mathbf{x}^* = H\mathbf{x}^* + \mathbf{v}$ and we are given the scheme

$$\mathbf{x}^{k+1} = H\mathbf{x}^k + \mathbf{v}, \qquad \mathbf{x}^0, \mathbf{v} \in \mathbb{R}^n. \tag{4}$$

Then $\mathbf{x}^k \to \mathbf{x}^*$ for any choice of $\mathbf{x}^0$ if and only if $\rho(H) < 1$.

Note: Of course, we would like to know not just convergence but the rate of it. For example, we achieve convergence with

$$H = \left[\begin{array}{cc} 0.99 & 10^6 \\ 0 & 0.99 \end{array} \right],$$

but it will take quite a long time. We will discuss this topic briefly later on.

Both of these methods are versions of splitting which can be applied to any A with nonzero diagonal elements. We write A as the sum of three matrices $L_0 + D + U_0$: subdiagonal (strictly lower-triangular), diagonal and superdiagonal (strictly upper-triangular) portions of A, respectively.

The Jacobi method

1) *Jacobi method.* We set $A - B = D$, the diagonal part of A, and we obtain the next iteration by solving the diagonal system

$$D\mathbf{x}^{(k+1)} = -(L_0 + U_0)\mathbf{x}^{(k)} + \mathbf{b}, \qquad H_{\mathrm{J}} = -D^{-1}(L_0 + U_0).$$

The Gauss–Seidel method

2) *Gauss–Seidel method*. We take $A - B = L_0 + D = L$, the lower-triangular part of A, and we generate the sequence $(\boldsymbol{x}^{(k)})$ by solving the triangular system

$$(L_0 + D)\,\boldsymbol{x}^{(k+1)} = -U_0\boldsymbol{x}^{(k)} + \boldsymbol{b}, \qquad H_{\mathrm{GS}} = -(L_0 + D)^{-1}U_0\,.$$

There is no need to invert $(L_0 + D)$, we calculate the components of $\boldsymbol{x}^{(k+1)}$ in sequence by forward substitution:

$$a_{ii}x_i^{(k+1)} = -\sum_{j<i} a_{ij}x_j^{(k+1)} - \sum_{j>i} a_{ij}x_j^{(k)} + b_i, \qquad i = 1..n.$$

Convergence

As we mentioned above, the sequence $x^{(k)}$ converges to the solution of $Ax = b$ if the spectral radius of the iteration matrix,

$$H_{\mathrm{J}} = -D^{-1}(L_0 + U_0) \text{ or } H_{\mathrm{GS}} = -(L_0 + D)^{-1}U_0,$$

respectively, is less than one. Our next goal is to prove that this is the case for two important classes of matrices A:

   a) diagonally dominant   and   b) positive definite matrices.

We start with recalling the simple, but very useful Gershgorin theorem.

Revision – Gershgorin theorem

All eigenvalues of an $n \times n$ matrix A are contained in the union of the Gershgorin discs in the complex plane:

$$\sigma(A) \subset \cup_{i=1}^{n} \Gamma_i, \quad \Gamma_i := \{z \in \mathbb{C} : |z - a_{ii}| \leq r_i\}, \quad r_i := \sum_{j \neq i} |a_{ij}|.$$

Strictly diagonally dominant matrices

Definition 4 (Strictly diagonally dominant matrices)

A matrix A is called strictly diagonally dominant by rows (resp. by columns) if

$$|a_{ii}| > \sum_{j \neq i} |a_{ij}|, \quad i = 1..n \qquad (\text{resp.} \quad |a_{jj}| > \sum_{i \neq j} |a_{ij}|, \quad j = 1..n).$$

From Gershgorin theorem, it follows that strictly diagonally dominant matrices are nonsingular.

Convergence of iterations

Theorem 5

If A is strictly diagonally dominant, then both the Jacobi and the Gauss-Seidel methods converge.

Proof. For the Gauss-Seidel method, the eigenvalues of the iteration matrix $H_{\mathrm{GS}} = -(L_0 + D)^{-1}U_0$ satisfy the equation

$$\det[H_{\mathrm{GS}} - \lambda I] = \det[-(L_0 + D)^{-1}U_0 - \lambda I] = 0.$$

Moreover,

$$\det[-(L_0+D)^{-1}U_0-\lambda I] = 0 \quad \Rightarrow \quad \det[A_\lambda] := \det[U_0+\lambda D+\lambda L_0] = 0.$$

It is easy to see that if $A = L_0 + D + U_0$ is strictly diagonally dominant, then for $|\lambda| \geq 1$ the matrix $A_\lambda = \lambda L_0 + \lambda D + U_0$ is strictly diagonally dominant too, hence it is nonsingular, and therefore the equality $\det[A_\lambda] = 0$ is impossible. Thus $|\lambda| < 1$, hence convergence. The proof for the Jacobi method is the same. $\square$

The Householder–John theorem

Theorem 6 (The Householder–John theorem)

If A and B are real matrices such that both A and $A - B - B^T$ are symmetric positive definite, then the spectral radius of $H = -(A - B)^{-1}B$ is strictly less than one.

Proof. Let λ be an eigenvalue of H, so $H\mathbf{w} = \lambda\mathbf{w}$ holds, where $\mathbf{w} \neq 0$ is an eigenvector. (Note that both λ and $\mathbf{w}$ may have nonzero imaginary parts when H is not symmetric, e.g. in the Gauss–Seidel method.) The definition of H provides equality $-B\mathbf{w} = \lambda(A - B)\mathbf{w}$, and we note that $\lambda \neq 1$ since otherwise A would be singular (which it is not). Thus, we deduce

$$\overline{\mathbf{w}}^T B \mathbf{w} = \frac{\lambda}{\lambda - 1}\overline{\mathbf{w}}^T A \mathbf{w}, \qquad (5)$$

where the bar means complex conjugation.

The Householder–John theorem

Proof. Cont. Moreover, writing $\boldsymbol{w} = \boldsymbol{u} + i\boldsymbol{v}$, where $\boldsymbol{u}$ and $\boldsymbol{v}$ are real, we find (for $C = C^T$) the identity $\overline{\boldsymbol{w}}^T C \boldsymbol{w} = \boldsymbol{u}^T C \boldsymbol{u} + \boldsymbol{v}^T C \boldsymbol{v}$, so symmetric positive definiteness in the assumption implies $\overline{\boldsymbol{w}}^T A \boldsymbol{w} > 0$ and $\overline{\boldsymbol{w}}^T (A - B - B^T) \boldsymbol{w} > 0$. In the latter inequality, we use relation (5) and its conjugate transpose to obtain

$$
0 < \overline{\boldsymbol{w}}^T A \boldsymbol{w} - \overline{\boldsymbol{w}}^T B \boldsymbol{w} - \overline{\boldsymbol{w}}^T B^T \boldsymbol{w} = \left(1 - \frac{\lambda}{\lambda - 1} - \frac{\overline{\lambda}}{\overline{\lambda} - 1} \right) \overline{\boldsymbol{w}}^T A \boldsymbol{w}
$$
$$
= \frac{1 - |\lambda|^2}{|\lambda - 1|^2} \, \overline{\boldsymbol{w}}^T A \boldsymbol{w}.
$$

Now $\lambda \neq 1$ implies $|\lambda - 1|^2 > 0$. Hence, recalling that $\overline{\boldsymbol{w}}^T A \boldsymbol{w} > 0$, we see that $1 - |\lambda|^2$ is positive. Therefore $|\lambda| < 1$ occurs for every eigenvalue of H as required. $\qquad\square$

The Householder–John theorem – A corollary

Corollary 7

1) *If A is symmetric positive definite, then the Gauss-Seidel method converges.*

2) *If both A and $2D - A$ are symmetric positive definite, then the Jacobi method converges.*

Proof. 1) For the Gauss-Seidel method, we take $A - B = L_0 + D$, thus $B = U_0$ is the superdiagonal part of symmetric A, hence $A - B - B^T$ is equal to D, the diagonal part of A, and if A is positive definite, then D is positive definite too (this is the first part of the Exercise 23 from Example Sheets).

2) For the Jacobi method, we have $B = A - D$, and if A is symmetric, then $A - B - B^T = 2D - A$. (The latter matrix is the same as A except that the signs of the off-diagonal elements are reversed.) $\square$

Linear systems in elliptic PDEs

As we have seen in the previous sections linear systems $A\boldsymbol{x} = \boldsymbol{b}$, where A is a real symmetric positive (negative) definite matrix, frequently occur in numerical methods for solving elliptic partial differential equations.

Poisson's equation on a square

A typical example we already encountered is Poisson's equation on a square where the *five-point formula* approximation yields an $n \times n$ system of linear equations with $n = m^2$ unknowns $u_{p,q}$:

$$u_{p-1,q} + u_{p+1,q} + u_{p,q-1} + u_{p,q+1} - 4u_{p,q} = h^2 f(ph, qh) \qquad (6)$$

In the *natural ordering*, when the grid points are arranged by columns, A is the following block tridiagonal matrix:

$$A = \begin{bmatrix} B & I & & & \\ I & B & I & & \\ & \ddots & \ddots & \ddots & \\ & & I & B & I \\ & & & I & B \end{bmatrix}, \qquad B = \begin{bmatrix} -4 & 1 & & & \\ 1 & -4 & 1 & & \\ & \ddots & \ddots & \ddots & \\ & & 1 & -4 & 1 \\ & & & 1 & -4 \end{bmatrix}. \qquad (7)$$

Lemma 8
For any ordering of the grid points, the matrix A of the system (6) is symmetric and negative definite.

Poisson's equation on a square

Note that when p or q is equal to 1 or m, then the values $u_{0,q}$, $u_{p,0}$ or $u_{p,m+1}$, $u_{m+1,q}$ are known boundary values and they should be moved to the right-hand side, thus leaving fewer unknowns on the left.

For any ordering of the grid points (ph, qh) we have shown in Lemma 8 that the matrix A of this linear system is symmetric and negative definite.

Corollary 9

For the linear system (6), for any ordering of the grid, both Jacobi and Gauss-Seidel methods converge.

Proof. By Lemma 8, A is symmetric and negative definite, hence convergence of Gauss-Seidel. To prove convergence of the Jacobi method, we need negative definiteness of the matrix $2D - A$, and that follows by the same arguments as in Lemma 8: recall that the proof operates with the modulus of the off-diagonal elements and does not depend on their sign. $\qquad\square$

Relaxation

It is often possible to improve the efficiency of the splitting method by *relaxation*. Specifically, instead of letting
$(A - B)\boldsymbol{x}^{(k+1)} = -B\boldsymbol{x}^{(k)} + \boldsymbol{b}$, we let

$$(A-B)\widehat{\boldsymbol{x}}^{(k+1)} = -B\boldsymbol{x}^{(k)}+\boldsymbol{b}, \quad \text{and then} \quad \boldsymbol{x}^{(k+1)} = \omega\widehat{\boldsymbol{x}}^{(k+1)}+(1-\omega)\boldsymbol{x}^{(k)}$$

where ω is a real constant called the *relaxation parameter*.

Note that $\omega = 1$ corresponds to the standard "unrelaxed" iteration. Good choice of ω leads to a smaller spectral radius of the iteration matrix (compared with the "unrelaxed" method), and the smaller the spectral radius, the faster the iteration converges.

Relaxation

To this end, let us express the relaxation iteration matrix H_ω in terms of $H = -(A - B)^{-1}B$. We have

$$\widehat{x}^{(k+1)} = Hx^{(k)} + v \quad \Rightarrow \quad \begin{aligned} x^{(k+1)} &= \omega\widehat{x}^{(k+1)} + (1 - \omega)x^{(k)} \\ &= \omega Hx^{(k)} + (1 - \omega)x^{(k)} + \omega v, \end{aligned}$$

hence

$$H_\omega = \omega H + (1 - \omega)I.$$

It follows that the spectra of H_ω and H are related by the rule $\lambda_\omega = \omega\lambda + (1 - \omega)$,

https://users.math.msu.edu/users/maccluer/Class/415/
SpectralMapping.pdf

therefore one may try to choose $\omega \in \mathbb{R}$ to minimize

$$\rho(H_\omega) = \max\left\{|\omega\lambda + (1 - \omega)| : \lambda \in \sigma(H)\right\}.$$

In general, $\sigma(H)$ is unknown, but often we have some information about it which can be utilized to find a "good" (rather than "best") value of ω. For example, suppose that it is known that $\sigma(H)$ is real and resides in the interval $[\alpha, \beta]$ where $-1 < \alpha < \beta < 1$. In that case we seek ω to minimize

$$\max \left\{ |\omega\lambda + (1 - \omega)| : \lambda \in [\alpha, \beta] \right\}.$$

It is readily seen that, for a fixed $\lambda < 1$, the function $f(\omega) = \omega\lambda + (1-\omega)$ is decreasing, therefore, as ω increases (decreases) from 1 the spectrum of H_ω moves to the left (to the right) of the spectrum of H. It is clear that the optimal location of the spectrum $\sigma(H_\omega)$ (or of the interval $[\alpha_\omega, \beta_\omega]$ that contains $\sigma(H_\omega)$) is the one which is centralized around the origin:

$$-[\omega\alpha + (1-\omega)] = \omega\beta + (1-\omega) \quad \Rightarrow \quad \begin{aligned} \omega_{\mathrm{opt}} &= \tfrac{2}{2-(\alpha+\beta)} \\ -\alpha_{\omega_{\mathrm{opt}}} = \beta_{\omega_{\mathrm{opt}}} &= \tfrac{\beta-\alpha}{2-(\alpha+\beta)} . \end{aligned}$$

Minimization of quadratic function

The methods we considered so far for solving $A\boldsymbol{x} = \boldsymbol{b}$, namely Jacobi, Gauss-Seidel, and those with relaxation, fit into the scheme

$$\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} + c_k \boldsymbol{d}^{(k)},$$

where we were aimed at getting $\rho(H) < 1$ for the iteration matix H. Say, for Jacobi with relaxation, we set $c_k = \omega$ and $\boldsymbol{d}^{(k)} = D^{-1}(\boldsymbol{b} - A\boldsymbol{x}^{(k)})$.

For solving $A\boldsymbol{x} = \boldsymbol{b}$ with a (positive definite) matrix $A > 0$, there is a different approach to constructing good iterative methods. It is based on succesive minimization of the quadratic function

$$F(\boldsymbol{x}^{(k)}) := \|\boldsymbol{x}^* - \boldsymbol{x}^{(k)}\|_A^2 = \|\boldsymbol{e}^{(k)}\|_A^2,$$

since the minimizer is clearly the exact solution. Here, $\|\boldsymbol{y}\|_A := (A\boldsymbol{y}, \boldsymbol{y})^{1/2} := \sqrt{\boldsymbol{y}^T A \boldsymbol{y}}$ is a Euclidean-type distance which is well-defined for $A > 0$.

Minimization of quadratic function

So, at each step k, we are decreasing the A-distance between $\boldsymbol{x}^{(k)}$ and the exact solution $\boldsymbol{x}^*$. Thus, for a symmetric positive definite $A > 0$, we choose an iterative method that provides the descent condition

$$\boldsymbol{x}^{(k+1)} = \boldsymbol{x}^{(k)} + c_k \boldsymbol{d}^{(k)} \quad \Rightarrow \quad F(\boldsymbol{x}^{(k+1)}) < F(\boldsymbol{x}^{(k)}). \qquad (8)$$

Minimization of quadratic function

An equivalent approach is to minimize the quadratic function

$$F_1(\boldsymbol{x}) = \tfrac{1}{2}\boldsymbol{x}^T A\boldsymbol{x} - \boldsymbol{x}^T \boldsymbol{b},$$

which attains its minimum when $\nabla F_1(\boldsymbol{x}) = A\boldsymbol{x} - \boldsymbol{b} = 0$, and which does not involve the unknown $\boldsymbol{x}^*$. It is easy to check that $F_1(\boldsymbol{x}) = \tfrac{1}{2}F(\boldsymbol{x}) - \tfrac{1}{2}c$, where $c = \boldsymbol{x}^{*T} A\boldsymbol{x}^*$ is a constant independent of k, hence equivalence.

Quadratic function – Jacobi and Gauss–Seidel

Both the Jacobi and the Gauss–Seidel methods satisfy (8), precisely

$$(A\boldsymbol{e}^{(k+1)}, \boldsymbol{e}^{(k+1)}) = (A\boldsymbol{e}^{(k)}, \boldsymbol{e}^{(k)}) - (C\boldsymbol{y}^{(k)}, \boldsymbol{y}^{(k)}) < (A\boldsymbol{e}^{(k)}, \boldsymbol{e}^{(k)}),$$

where for Gauss-Seidel: $\quad C = D > 0, \qquad \boldsymbol{y}^{(k)} := (L_0 + D)^{-1} A \boldsymbol{e}^{(k)};$

and for Jacobi: $\quad C = 2D - A > 0, \quad \boldsymbol{y}^{(k)} := D^{-1} A \boldsymbol{e}^{(k)}.$

A-**orthogonal projection method:** Next, we strengthen the descent condition (8), namely given $\mathbf{x}^{(k)}$ and some $\mathbf{d}^{(k)}$ (called a *search direction*), we will seek $\mathbf{x}^{(k+1)}$ from the set of vectors on the line $\ell = \{\mathbf{x}^{(k)} + \alpha \mathbf{d}^{(k)}\}_{\alpha \in \mathbb{R}}$ such that it makes the value of $F(\mathbf{x}^{(k+1)})$ not just smaller than $F(\mathbf{x}^{(k)})$, but as small as possible (with respect to this set), namely

$$\mathbf{x}^{(k+1)} := \arg \min_{\alpha} F(\mathbf{x}^{(k)} + \alpha \mathbf{d}^{(k)}). \tag{9}$$

Lemma 10
The minimizer in (9) is given by the formula

$$\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}, \qquad \alpha_k = \frac{(\mathbf{r}^{(k)}, \mathbf{d}^{(k)})}{(A\mathbf{d}^{(k)}, \mathbf{d}^{(k)})}. \qquad (10)$$

A-orthogonal projection

Proof. From the definition of F, it follows that in (9) we should choose the point $\boldsymbol{x}^{(k+1)} \in \ell$ that minimizes the A-distance between $\boldsymbol{x}^*$ and the points $\boldsymbol{y} \in \ell$. Geometrically, it is clear that the minimum occurs when $\boldsymbol{x}^{(k+1)}$ is the A-orthogonal projection of $\boldsymbol{x}^*$ onto the line $\ell = \{\boldsymbol{x}^{(k)} + \alpha \boldsymbol{d}^{(k)}\}$, i.e., when

$$
\begin{aligned}
\boldsymbol{x}^* - \boldsymbol{x}^{(k+1)} \perp_A \boldsymbol{d}^{(k)} \quad &\Rightarrow \quad A(\boldsymbol{x}^* - \boldsymbol{x}^{(k+1)}) \perp \boldsymbol{d}^{(k)} \\
&\Rightarrow \quad \boldsymbol{r}^{(k+1)} = \boldsymbol{r}^{(k)} - \alpha_k A \boldsymbol{d}^{(k)} \perp \boldsymbol{d}^{(k)} .
\end{aligned}
$$

This gives expression for α_k in (10). $\qquad\square$