Numerical Analysis - Part II

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Lecture 24

Multigrid methods

The speed of convergence of some iterative methods (Jacobi with relaxation, Gauss-Seidel, etc.) can be increased drastically when the linear system originates in the discretization of PDEs, using *multigrid methods*. Here we look at the system $A\mathbf{u} = \mathbf{b}$ originating from the 3-point formula for the Poisson equation on an *m*-grid $\Omega_h = \{ih : 1 \le i \le m\}$, h = 1/(m+1), being solved by the weighted Jacobi iteration. Recall that the matrix A in this case is given by

$$A = \begin{bmatrix} 2 & -1 & & \\ -1 & 2 & -1 & & \\ & \ddots & & \\ & & -1 & 2 & -1 \\ & & & -1 & 2 \end{bmatrix} \in \mathbb{R}^{m \times m}.$$

The diagonal part of A is D = 2I. Thus the weighted Jacobi iterations takes the form:

$$\boldsymbol{u}^{(\nu+1)} = H_{\omega} \boldsymbol{u}^{(\nu)} + (\omega/2) \boldsymbol{b}$$

where $\nu = 0, 1, ...$ is the iteration count, $H = I - D^{-1}A = I - \frac{1}{2}A$, and $H_{\omega} = \omega H + (1 - \omega)I = I - \frac{\omega}{2}A$. The error decay is expressed in terms of the iteration matrix H_{ω} :

$$\boldsymbol{e}^{(
u)}=H^{
u}_{\omega}\boldsymbol{e}^{(0)}.$$

We know from the results of Lecture 2 that the eigenvectors and the eigenvalues of H_{ω} are

$$\boldsymbol{w}^{k} = \left[\sin i \frac{k\pi}{m+1}\right]_{i=1,\dots,m}, \qquad \lambda_{k}(\omega) = 1 - 2\omega \sin^{2} \frac{k\pi}{2(m+1)} \qquad (k = 1,\dots,m).$$

Multigrid methods – Eigenvectors and eigenvalues

Consider the choice $\omega = 1/2$; then the eigenvalues of H_{ω} are $\lambda_k = 1 - \sin^2 \frac{k\pi}{2(m+1)} = \cos^2 \frac{k\pi}{2(m+1)}$. With this choice, the eigenvalues are all positive and decreasing with k, see Figure below.

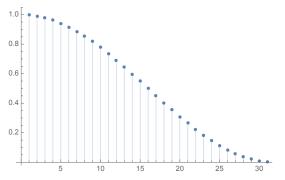


Figure: Eigenvalues of H_{ω} for $\omega = 1/2$ (m = 31).

In particular $\rho(H_{\omega}) = \lambda_1 = \cos^2 \frac{\pi}{2(m+1)} \approx 1 - \frac{\pi^2}{4m^2} < 1$, guaranteeing convergence, although a very slow one when *m* is large!

However, expanding the error with respect to the (orthogonal) eigenvectors we obtain

$$\boldsymbol{e}^{(\nu)} = \sum_{k=1}^{m} a_{k}^{(\nu)} \boldsymbol{w}^{k}, \qquad \boldsymbol{e}^{(\nu)} = H_{\omega}^{\nu} \boldsymbol{e}^{(0)} \quad \Rightarrow \quad |a_{k}^{(\nu)}| = |\lambda_{k}|^{\nu} |a_{k}^{(0)}|,$$

i.e. the components of $e^{(\nu)}$ (with respect to the basis of eigenvectors) decay at a different rate for different frequencies $k = 1, \ldots, m$. More precisely, the high frequencies, where k is close to m, will decay faster than the low frequencies, where k is closer to 1.

Let us say that $k \in (0, m + 1) = (0, \frac{1}{h})$ is high frequency (HF) with respect to the grid Ω_h if $kh \ge 1/2$ (i.e., $\frac{m+1}{2} \le k \le m$). Then the decay rate for the high frequency components of the error e is at least:

$$\mu_* = |\lambda_{(m+1)/2}| = 1 - \sin^2(\pi/4) = 1/2.$$

Therefore, for the coefficients at the HF components of $e^{(\nu)}$ we obtain

$$|a_k^{(
\nu)}| \le |\mu_*|^{
u} |a_k^{(0)}| = \left(rac{1}{2}
ight)^{
u} |a_k^{(0)}| \ll |a_k^{(0)}| \,,$$

i.e. the Jacobi method converges fast for high frequencies.

The main observation of the multigrid is to note that the low frequencies $k \in (\frac{1}{4h}, \frac{1}{2h})$ with respect to the grid Ω_h become high frequencies for the *coarser grid* Ω_{2h} with step 2h; indeed for such k we have $k(2h) \ge 1/2$.

Changing the grid

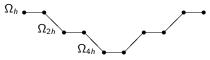
The idea of the multigrid method then is that, although the global error may decrease slowly by iteration, its components with high frequencies relative to Ω_h are suppressed very quickly, and that dealing with the remaining components (with low frequencies relative to Ω_h) we can move to the coarse grid Ω_{2h} , where these components (in part) would be of high frequencies, and thus they can be suppressed in a similar way. Therefore, we cover the domain [0, 1] by a range of nested grids, of increasing coarseness, say,

$$\Omega_h \subset \Omega_{2h} \subset \Omega_{4h} \subset \cdots \subset \Omega_{2^j h}.$$

At every Ω_{h_i} , the iterations (Jacobi, or Gauss-Seidel) remove the high frequencies relative to this grid, and we move to Ω_{2h_i} . On the coarsest grid, where the number of variables is small, we can afford to solve the equations with a direct method, by Cholesky, say.

Remark 1

If we follow the recursive procedure outlined below, then we see that the algorithm starts at the finest grid, travels to the coarsest (where we apply a direct solver), and back to the finest:



For this reason, the algorithm above describes what is known as a V-cycle, hence the name of the routine.

Multigrid methods – The procedure

A typical multigrid method can be summarized by the following routine **MGV**, which gives an approximate solution to the linear system Ax = b, starting from the initial guess u^0 . We assume below that the size of the linear system is $m = 1/h - 1 = 2^{\ell} - 1$ for some integer ℓ .

 $\mathsf{MGV}(A,b,\pmb{u}^0)$

- 1. If size of A is small enough, use a direct method to solve Ax = b and exit. Else:
- 2. Presmoothing: Perform a small number (typically \leq 5) of Jacobi or Gauss-Seidel iterations on Au = b starting from u^0 .
- 3. Let $\mathbf{r} = \mathbf{b} A\mathbf{u}$ be the residual, with \mathbf{u} from the previous step.
- 4. Let $I_{2h}^h : \mathbb{R}^{\frac{m+1}{2}-1} \to \mathbb{R}^m$ be an *interpolation operator* that interpolates vectors on the coarse grid Ω_{2h} to vectors on the fine grid Ω_h ; and let $R_{h}^{2h} : \mathbb{R}^m \to \mathbb{R}^{(m+1)/2-1}$ be a *restriction operator* that restricts vectors on the fine grid Ω_h to vectors on the coarse grid Ω_{2h} .
- 5. Let $\tilde{A} = R_h^{2h} A I_{2h}^h$ which is of size $\approx m/2 \times m/2$.
- 6. Let $\tilde{\delta} = MGV(\tilde{A}, R_h^{2h} r, 0)$ (approximate solution to the residual equation $A\delta = r$ on the coarse grid)
- 7. Let $\boldsymbol{u} = \boldsymbol{u} + l_{2h}^h \tilde{\delta}$
- 8. Postsmoothing: apply a few Jacobi or Gauss-Seidel iterations starting from \boldsymbol{u} on $Ax = \boldsymbol{b}$
- 9. Return u

https://en.wikipedia.org/wiki/Multigrid_method

Interpolation

To make the algorithm above complete, one needs to define the interpolation and restriction operators. A common choice for the interpolation operator is linear interpolation, i.e.,

$$(I_{2h}^h \mathbf{v})_{2i} = \mathbf{v}_i$$
 and $(I_{2h}^h \mathbf{v})_{2i+1} = (\mathbf{v}_i + \mathbf{v}_{i+1})/2$

for $i \in [0, \frac{1}{2h})$. Assuming zero boundary conditions, this takes the following matrix form:

	1]	
$I_{2h}^{h} = 1/2$	2				
	1	1			
		2			
		1	1		
			÷		
			•		
			1	1	
				2	
	L			1	

Restriction

For the restriction operator R_h^{2h} , we take the following averaging operator

$$(R_h^{2h}\mathbf{v})_i = \frac{1}{4}(\mathbf{v}_{2i-1} + 2\mathbf{v}_{2i} + \mathbf{v}_{2i+1}) \qquad i \in (0, \frac{1}{2h})$$

which corresponds, up to scaling, to the transpose of the linear interpolation operator above; namely $R_h^{2h} = \frac{1}{2} (I_{2h}^h)^T$. In practice, one may need to apply the routine **MGV** multiple times to improve the accuracy, each time starting from the previously obtained solution.

https://www.damtp.cam.ac.uk/user/hf323/M21-II-NA/ demos/multigrid/multigrid.html