

Mathematical Tripos Part II: Michaelmas Term 2024

Numerical Analysis – Lecture 19

Algorithm 4.26 (The conjugate gradient method) Here it is.

(A) For any initial vector $\mathbf{x}^{(0)}$, set $\mathbf{d}^{(0)} = \mathbf{r}^{(0)} = \mathbf{b} - A\mathbf{x}^{(0)}$;

(B) For $k \geq 0$, calculate $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}$ and the residual

$$\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k A\mathbf{d}^{(k)}, \quad \text{with} \quad \alpha_k := \{\mathbf{r}^{(k+1)} \perp \mathbf{d}^{(k)}\} = \frac{(\mathbf{r}^{(k)}, \mathbf{d}^{(k)})}{(A\mathbf{d}^{(k)}, \mathbf{d}^{(k)})}, \quad k \geq 0. \quad (4.8)$$

(C) For the same k , the next conjugate direction is the vector

$$\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)}, \quad \text{with} \quad \beta_k := \{\mathbf{d}^{(k+1)} \perp A\mathbf{d}^{(k)}\} = -\frac{(\mathbf{r}^{(k+1)}, A\mathbf{d}^{(k)})}{(\mathbf{d}^{(k)}, A\mathbf{d}^{(k)})}, \quad k \geq 0. \quad (4.9)$$

Theorem 4.27 (Properties of CGM) For every $m \geq 0$, the conjugate gradient method has the following properties.

(1) The linear space spanned by the residuals $\{\mathbf{r}^{(i)}\}$ is the same as the linear space spanned by the conjugate directions $\{\mathbf{d}^{(i)}\}$ and it coincides with the space spanned by $\{A^i \mathbf{r}^{(0)}\}$:

$$\text{span}\{\mathbf{r}^{(i)}\}_{i=0}^m = \text{span}\{\mathbf{d}^{(i)}\}_{i=0}^m = \text{span}\{A^i \mathbf{r}^{(0)}\}_{i=0}^m.$$

(2) The residuals satisfy the orthogonality conditions: $(\mathbf{r}^{(m)}, \mathbf{r}^{(i)}) = (\mathbf{r}^{(m)}, \mathbf{d}^{(i)}) = 0$ for $i < m$.

(3) The directions are conjugate (A -orthogonal): $(\mathbf{d}^{(m)}, \mathbf{d}^{(i)})_A = (\mathbf{d}^{(m)}, A\mathbf{d}^{(i)}) = 0$ for $i < m$.

Proof. We use induction on $m \geq 0$, the assertions being trivial for $m = 0$, since $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$ and (2)-(3) are void. Therefore, assuming that the assertions are true for some $m = k$, we ask if they remain true when $m = k + 1$.

(1) Formula (4.9)

$$\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)}$$

readily implies that equivalence of the spaces spanned by $(\mathbf{r}^{(i)})_0^k$ and $(\mathbf{d}^{(i)})_0^k$, is preserved when k is increased to $k + 1$. Similarly, from $\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k A\mathbf{d}^{(k)}$ in (4.8), and from the inductive assumption $\mathbf{r}^{(k)}, \mathbf{d}^{(k)} \in \text{span}\{A^i \mathbf{r}^{(0)}\}_{i=0}^k$, it follows that $\mathbf{r}^{(k+1)} \in \text{span}\{A^i \mathbf{r}^{(0)}\}_{i=0}^{k+1}$.

(2) Turning to assertion (2), we need $\mathbf{r}^{(k+1)} \perp \mathbf{r}^{(i)}$ for $i \leq k$, which by (1) is equivalent to

$$\mathbf{r}^{(k+1)} \perp \mathbf{d}^{(i)} \quad \text{for} \quad i \leq k.$$

We have $\mathbf{r}^{(k+1)} \perp \mathbf{d}^{(k)}$ by the definition of α_k in (4.8), so we need

$$\mathbf{r}^{(k+1)} \stackrel{(4.8)}{=} \mathbf{r}^{(k)} - \alpha_k A\mathbf{d}^{(k)} \perp \mathbf{d}^{(i)} \quad \text{for} \quad i < k,$$

and this follow from the inductive assumptions $\mathbf{r}^{(k)} \perp \mathbf{d}^{(i)}$ and $A\mathbf{d}^{(k)} \perp \mathbf{d}^{(i)}$.

(3) It remains to justify (3), namely that $\mathbf{d}^{(k+1)}$ defined in (4.9) satisfies

$$\mathbf{d}^{(k+1)} \perp A\mathbf{d}^{(i)} \quad \text{for} \quad i \leq k.$$

The value of β_k in (4.9) is defined to give $\mathbf{d}^{(k+1)} \perp A\mathbf{d}^{(k)}$, so we need

$$\mathbf{d}^{(k+1)} \stackrel{(4.9)}{=} \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)} \perp A\mathbf{d}^{(i)} \quad \text{for} \quad i < k.$$

By the inductive hypothesis $\mathbf{d}^{(k)} \perp A\mathbf{d}^{(i)}$, hence it remains to establish that $\mathbf{r}^{(k+1)} \perp A\mathbf{d}^{(i)}$ for $i < k$. Now, the formula (4.8) yields $A\mathbf{d}^{(i)} = (\mathbf{r}^{(i)} - \mathbf{r}^{(i+1)})/\alpha_i$, therefore we require the conditions $\mathbf{r}^{(k+1)} \perp (\mathbf{r}^{(i)} - \mathbf{r}^{(i+1)})$ for $i < k$, and they are a consequence of the assertion (2) for $m = k + 1$ obtained previously. \square

Corollary 4.28 (A termination property) *If the conjugate gradient method is applied in exact arithmetic, then, for any $\mathbf{x}^{(0)} \in \mathbb{R}^n$, termination occurs after at most n iterations. More precisely, termination occurs after at most s iterations, where $s = \dim \text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$ (which can be smaller than n).*

Proof. Assertion (2) of Theorem 4.27 states that residuals $(\mathbf{r}^{(k)})_{k \geq 0}$ form a sequence of mutually orthogonal vectors in \mathbb{R}^n , therefore at most n of them can be nonzero. Since they also belong to the space $\text{span}\{A^i \mathbf{r}_0\}_{i=0}^{n-1}$, their number is bounded by the dimension of that space. \square

Definition 4.29 (The Krylov subspaces) Let A be an $n \times n$ matrix, $\mathbf{v} \in \mathbb{R}^n$ nonzero, and $m \in \mathbb{N}$. The linear space $K_m(A, \mathbf{v}) := \text{span}\{A^i \mathbf{v}\}_{i=0}^{m-1}$ is called the m -th Krylov subspace of \mathbb{R}^n .

Theorem 4.30 (Number of iterations in CGM) *Let $A > 0$, and let s be the number of its distinct eigenvalues. Then, for any \mathbf{v} ,*

$$\dim K_m(A, \mathbf{v}) \leq s \quad \forall m. \quad (4.10)$$

Hence, for any $A > 0$, the number of iterations of the CGM for solving $A\mathbf{x} = \mathbf{b}$ is bounded by the number of distinct eigenvalues of A .

Proof. Inequality (4.10) is true not just for positive definite $A > 0$, but for any A with n linearly independent eigenvectors (\mathbf{u}_i) . Indeed, in that case one can expand $\mathbf{v} = \sum_{i=1}^n a_i \mathbf{u}_i$, and then group together eigenvectors with the same eigenvalues: for each λ_ν we set $\mathbf{w}_\nu = \sum_{k=1}^{m_\nu} a_{i_k} \mathbf{u}_{i_k}$ if $A\mathbf{u}_{i_k} = \lambda_\nu \mathbf{u}_{i_k}$. Then

$$\mathbf{v} = \sum_{\nu=1}^s c_\nu \mathbf{w}_\nu, \quad c_\nu \in \{0, 1\},$$

hence $A^i \mathbf{v} = \sum_{\nu=1}^s c_\nu \lambda_\nu^i \mathbf{w}_\nu$, thus for any m we get $K_m(A, \mathbf{v}) \subseteq \text{span}\{\mathbf{w}_1, \mathbf{w}_2, \dots, \mathbf{w}_s\}$, and that proves (4.10). By Corollary 4.28, the number of iteration in CGM is bounded by $\dim K_m(A, \mathbf{r}^{(0)})$, hence the final conclusion. \square

Remark 4.31 Theorem 4.30 shows that, unlike other iterative schemes, the conjugate gradient method is both iterative and direct: each iteration produces a reasonable approximation to the exact solution, and the exact solution itself will be recovered after n iterations at most.

We now simplify and reformulate Algorithm 4.26.

Firstly, we rewrite expressions for the parameters α_k and β_k in (4.8)-(4.9) as follows:

$$\begin{aligned} \alpha_k &= \frac{(\mathbf{r}^{(k)}, \mathbf{d}^{(k)})}{(\mathbf{d}^{(k)}, A\mathbf{d}^{(k)})} \stackrel{(c)}{=} \frac{\|\mathbf{r}^{(k)}\|^2}{(\mathbf{d}^{(k)}, A\mathbf{d}^{(k)})} > 0, \\ \beta_k &= -\frac{(\mathbf{r}^{(k+1)}, A\mathbf{d}^{(k)})}{(\mathbf{d}^{(k)}, A\mathbf{d}^{(k)})} \stackrel{(a)}{=} -\frac{(\mathbf{r}^{(k+1)}, \mathbf{r}^{(k+1)} - \mathbf{r}^{(k)})}{(\mathbf{d}^{(k)}, \mathbf{r}^{(k+1)} - \mathbf{r}^{(k)})} \stackrel{(b)}{=} \frac{\|\mathbf{r}^{(k+1)}\|^2}{(\mathbf{d}^{(k)}, \mathbf{r}^{(k)})} \stackrel{(c)}{=} \frac{\|\mathbf{r}^{(k+1)}\|^2}{\|\mathbf{r}^{(k)}\|^2} > 0. \end{aligned}$$

Here, for β , we used in (a) the fact that $A\mathbf{d}^{(k)}$ is a multiple of $\mathbf{r}^{(k+1)} - \mathbf{r}^{(k)}$ by (4.8), and in (b) orthogonality of $\mathbf{r}^{(k+1)}$ to both $\mathbf{r}^{(k)}$, $\mathbf{d}^{(k)}$ proved in Theorem 4.27(2). Then, for both β and α , we used in (c) the property $(\mathbf{d}^{(k)}, \mathbf{r}^{(k)}) = \|\mathbf{r}^{(k)}\|^2$ which follows from (4.9) with index $k+1$, taking in account orthogonality $\mathbf{r}^{(k+1)} \perp \mathbf{d}^{(k)}$.

Secondly, we let $\mathbf{x}^{(0)}$ be the zero vector.

Algorithm 4.32 (Standard form of the conjugate gradient method) Here it is.

- (1) Set $k = 0$, $\mathbf{x}^{(0)} = 0$, $\mathbf{r}^{(0)} = \mathbf{b}$, and $\mathbf{d}^{(0)} = \mathbf{r}^{(0)}$;
- (2) Calculate the matrix-vector product $\mathbf{v}^{(k)} = A\mathbf{d}^{(k)}$ and $\alpha_k = \|\mathbf{r}^{(k)}\|^2 / (\mathbf{d}^{(k)}, \mathbf{v}^{(k)}) > 0$;
- (3) Apply the formulae $\mathbf{x}^{(k+1)} = \mathbf{x}^{(k)} + \alpha_k \mathbf{d}^{(k)}$ and $\mathbf{r}^{(k+1)} = \mathbf{r}^{(k)} - \alpha_k \mathbf{v}^{(k)}$;
- (4) Stop if $\|\mathbf{r}^{(k+1)}\|$ is acceptably small;
- (5) Set $\mathbf{d}^{(k+1)} = \mathbf{r}^{(k+1)} + \beta_k \mathbf{d}^{(k)}$, where $\beta_k = \|\mathbf{r}^{(k+1)}\|^2 / \|\mathbf{r}^{(k)}\|^2 > 0$;
- (6) Increase $k \rightarrow k+1$ and go back to (2).

The total work is dominated by the number of iterations, multiplied by the time it takes to compute $\mathbf{v}^{(k)} = A\mathbf{d}^{(k)}$. Thus the conjugate gradient algorithm is highly suitable when most of the elements of A are zero, i.e. when A is *sparse*.