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in very **simple geometries**, Cartesian

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Sometimes **FAST** transform + less modes needed → competitive

Two idea - as in FE

Spectral representation

$$u(x, t) = \sum^N \hat{u}_n(t) \phi_n(x)$$

with amplitudes $u_n(t)$ and basis functions $\phi_n(x)$, e.g. Fourier

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Galerkin approximation “weighted residuals”. For PDE

$$A(u) = f$$

require **residue** to be orthogonal to each ϕ_m :

$$\langle A(u) - f, \phi_m \rangle = 0 \quad \text{for } m = 1, \dots, N$$

Local vs Global

E.g. for Fourier

$$u(x) = \int e^{ikx} \hat{u}(k) dk \quad \hat{u}(k) = \frac{1}{2\pi} \int e^{-ikx} u(x) dx$$

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Poisson problem

$$\frac{d^2 u}{dx^2} = \rho \quad \text{expensive global problem in real space}$$

$$-k^2 \hat{u} = \hat{\rho} \quad \text{local in Fourier space}$$

Local/Global continued

Nonlinear terms and spatially vary coefficients

$u(x)v(x)$ local in real space

$$\widehat{uv}(k) = \frac{1}{2\pi} \int_{l+m=k} \hat{u}(l)\hat{v}(m) \quad \text{global in Fourier}$$

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local = cheap global = expensive

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Navier-Stokes has both local & global in real or Fourier – need compromise

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combines Fourier and real space operations

Evaluate the nonlinear term in real space, and in Fourier space evaluate derivatives and invert the Poisson problem.

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Alternative method of satisfying PDE at **collocation points** rather than in Galerkin projection.

Choice of spectral basis function $\phi_n(x)$

1. complete
2. orthogonal for some weight w

$$\langle \phi_n \phi_m \rangle = \int \phi_n \phi_m w(x) dx = N_n \delta_{nm}$$

3. smooth
4. fast convergence
5. FAST transform
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Strongly recommend

- ▶ Fully periodic \rightarrow Fourier, $e^{in\theta}$
- ▶ Finite interval \rightarrow Chebyshev $T_n(\cos \theta) = \cos n\theta$

Chebyshev polynomials

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$$T_0(x) = 1, \quad T_1(x) = x, \quad T_2(x) = 2x^2 - 1$$
$$T_3(x) = 4x^3 - 3x, \quad T_4(x) = 8x^4 - 8x^2 + 1$$

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$$(1-x^2) T_n'' - x T_n' + n^2 T_n = 0$$

$$T_{n+1} = 2xT_n - T_{n-1}$$

$$2T_n = \frac{1}{n+1} T_{n+1}' - \frac{1}{n-1} T_{n-1}'$$

Fourier series

Fully periodic (really defined on a circle):

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– awkward $\frac{1}{2}a_0$ if use sines and cosines.

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E.g.

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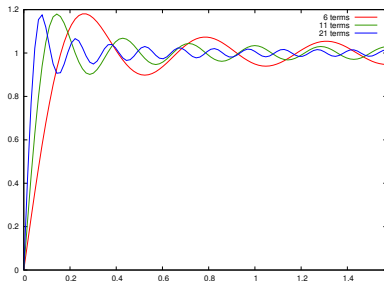
– convergence controlled by singularity of $f(\theta)$ in complex θ -plane

Gibbs phenomenon

Discontinuity \rightarrow poor $\sum \frac{\pm 1}{n}$ convergence

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with point-wise convergence

but 14% overshoot within $\frac{1}{N}$ of discontinuity

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Hence function $|x|$ on $-1 < x < 1$
becomes fully 2π periodic in $-\pi < \theta < 0$

Discrete Fourier Transform (DFT)

Odd $N = 2M + 1$.

Equi-spaced collocation points $\theta_j = \frac{2\pi j}{N}$ for $j = 1, \dots, N$

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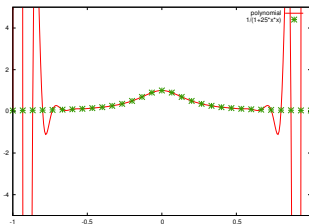
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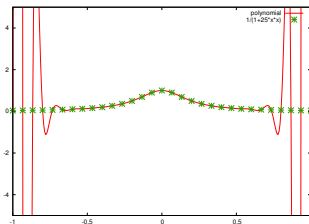
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Fitting polynomial through equi-spaced points can be **badly wrong** in between fitting points.



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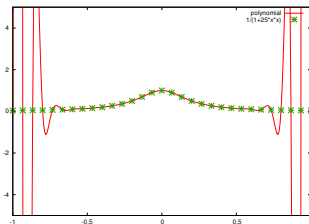
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However DFT well behaved, because effectively Chebyshev polynomials fitted at points $x_j = \cos(\pi j/N)$ – crowded at ends.

Aliasing

– counter rotating wagon wheels in strobe light

High $(N + k)$ frequency, e.g. $g(\theta) = e^{i(N+k)\theta}$,
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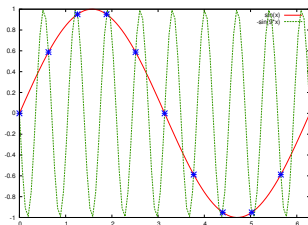
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E.g. $N = 10$ equispaced points cannot distinguish between $\sin \theta$
and $-\sin 9\theta$



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of exact Fourier modes \hat{f}_n in $n > M$

appear to DFT \tilde{f}_n

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In 3D throw away $\frac{19}{27}$ of the modes.

Fast Fourier Transform

DFT calculation for $n = -\frac{1}{2}N, \dots, \frac{1}{2}N$

$$\tilde{f}_n = \sum_{j=1}^N f(\theta_j) \omega^{nj}, \quad \text{with } \theta_j = \frac{2\pi j}{N} \text{ and } \omega = e^{i\theta_1}$$

looks like N coefficients \times sum of N terms $= N^2$ operations.

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$$= \sum_{k=1}^{N/2} f(\theta_{2k}) \omega_2^{nk} + \omega^{-1} \sum_{k=1}^{N/2} f(\theta_{2k-1}) \omega_2^{nk} \quad \text{with } \omega_2 = \omega^2$$

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which is 2 lots of DFT on $\frac{1}{2}N$ points $2(\frac{1}{2}N)^2 = \frac{1}{2}N^2$ operations

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which is 2 lots of DFT on $\frac{1}{2}N$ points $2(\frac{1}{2}N)^2 = \frac{1}{2}N^2$ operations

If $N = 2^K$, can half K times $\rightarrow N \ln_2 N$ operations.

Fast Fourier Transform

DFT calculation for $n = -\frac{1}{2}N, \dots, \frac{1}{2}N$

$$\tilde{f}_n = \sum_{j=1}^N f(\theta_j) \omega^{nj}, \quad \text{with } \theta_j = \frac{2\pi j}{N} \text{ and } \omega = e^{i\theta_1}$$

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Program: identify even/odd at each 2^n -level $n = 1, \dots, K$,
i.e. binary representation of j

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Also FFT speed up

Differential Matrix

To differentiate data with exponential accuracy

$$f(\theta_j) \xrightarrow{\text{transform}} \tilde{f}_n \xrightarrow{\text{differentiate}} n\tilde{f}_n \xrightarrow{\text{transform}} f'(\theta_j)$$

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NB $D^{(2)} \neq DD$

Navier-Stokes

$$\nabla \cdot \mathbf{u} = 0$$

$$\frac{\partial \mathbf{u}}{\partial t} + \mathbf{u} \cdot \nabla \mathbf{u} = -\nabla p + \nu \nabla^2 \mathbf{u}$$

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with $\widehat{\mathbf{u} \cdot \nabla \mathbf{u}}$ by pseudo-spectral real space evaluation

Boundary conditions

If homogeneous BCs, recombine to satisfy BCs

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$1/N^2$ crowding of $x_j = \cos \theta_j$ near ± 1

\rightarrow stability if $\Delta t < D/N^4$

Bridging the gap

Local

Finite Differences
point data

Finite Elements

FE h^p

Global

Spectral
whole interval

Bridging the gap

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Splines
global points

Wavelets
local waves