

Towards a classification theory for data-driven dynamical systems

Matthew Colbrook
University of Cambridge
28/10/2024

“To classify is to bring order into chaos.” - **George Pólya**

C., Mezić, Stepanenko *“Limits and Powers of Koopman Learning,”* arxiv preprint, 2024.

For papers and talk slides/videos, visit:
<http://www.damtp.cam.ac.uk/user/mjc249/home.html>

Data-driven dynamical systems

- Compact metric space (\mathcal{X}, d) – the state space

- $x \in \mathcal{X}$ – the state

cts $F: \mathcal{X} \rightarrow \mathcal{X}$ – the dynamics: $x_{n+1} = F(x_n)$

Dynamics (geometry)
19th century

- Borel measure ω on \mathcal{X}

- Function space $L^2 = L^2(\mathcal{X}, \omega)$ (elements g called “observables”)

- Koopman operator $\mathcal{K}_F: L^2 \rightarrow L^2$; $[\mathcal{K}_F g](x) = g(F(x))$

- **Available** snapshot data: $\left\{ \left(x^{(m)}, y^{(m)} = F(x^{(m)}) \right) : m = 1, \dots, M \right\}$

NB: Pointwise definition of \mathcal{K}_F needs $F\#\omega \ll \omega$ – this will hold throughout.

NB: \mathcal{K}_F bounded equivalent to $dF\#\omega/d\omega \in L^\infty$ – this will hold throughout (can be dropped).

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Analysis
20th century

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- Dynamics (geometry)
19th century
- Analysis
20th century
- Data
21st century

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Why you should care about Koopman

Fundamental in ergodic theory

Graduate Texts
in Mathematics

Peter Walters

An Introduction
to Ergodic Theory



$$x_{n+1} = F(x_n)$$

$$[\mathcal{K}g](x) = g(F(x))$$

E.g., key to ergodic theorems of Birkhoff and von Neumann.

Trades: Nonlinear, finite-dimensional \Rightarrow Linear, infinite-dimensional.

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Can provide a *diagonalization* of a nonlinear system.

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$$g(x) = \sum_{\text{eigenvalues } \lambda_j} c_{\lambda_j} \overset{\text{eigenfunction of } \mathcal{K}}{\varphi_{\lambda_j}(x)} + \int_{-\pi}^{\pi} \overset{\text{continuous spectrum}}{\phi_{\theta,g}(x)} d\theta$$

$$g(x_n) = [\mathcal{K}^n g](x_0)$$

$$= \sum_{\text{eigenvalues } \lambda_j} c_{\lambda_j} \lambda_j^n \varphi_{\lambda_j}(x_0) + \int_{-\pi}^{\pi} e^{in\theta} \phi_{\theta,g}(x_0) d\theta$$

Spectral properties encode: geometric features, invariant measures, transient behavior, long-time behavior, coherent structures, quasiperiodicity, etc.

Trades: Nonlinear, finite-dimensional \Rightarrow Linear, infinite-dimensional.

Why you should care about Koopman

Fundamental in ergodic theory

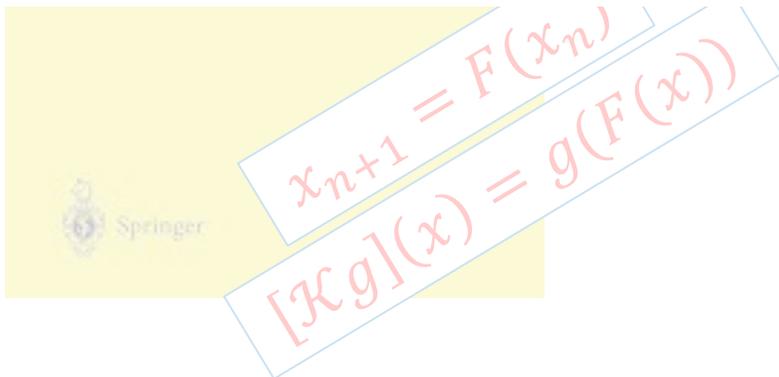
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+ HUGE recent interest in their spectral properties...



$$= \sum_{\text{eigenvalues } \lambda_j} c_{\lambda_j} \lambda_j^n \varphi_{\lambda_j}(x_0) + \int_{-\pi}^{\pi} e^{in\theta} \phi_{\theta,g}(x_0) d\theta$$

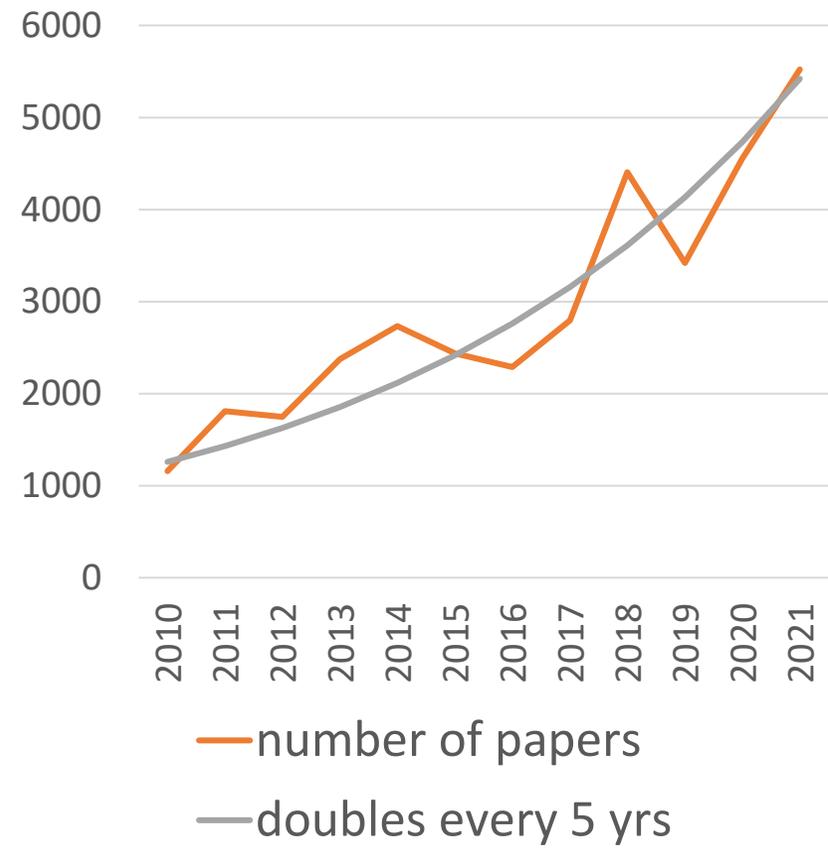
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Trades: Nonlinear, finite-dimensional \Rightarrow Linear, infinite-dimensional.

New Papers on "Koopman Operators"

Central in data-driven era



Loads of applications!!!

New Papers on "Koopman Operators"

Central in data-driven era



2012

Applied Koopmanism⁹⁾

Marko Budišić, Ryan Mohr, and Igor Mezic
 Department of Mechanical Engineering, University of California, Santa Barbara
 (Received 11 June 2012; accepted 30 November 2012; published online 11 December 2012)

A majority of methods from dynamical system analysis, especially Poincaré's geometric picture that focuses on "dynamics of stable manifolds," has shown difficulties in handling high-dimensional uncertain systems, which are more and more common in engineering "big data" measurements. This overview article presents an alternative picture, based on the "dynamics of observables" picture. This

CHAOS 22, 047510 (2012)



SIAM Review Vol. 64, No. 2, pp 229-340

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Modern Koopman Theory for

Issue: Many practitioners view it as a magic bullet, but standard algorithms typically fail to converge! (Inf-dim spectral problems.)

Question: When can we reliably learn Koopman spectral properties from system data, and when is it impossible?

—number of papers
 —doubles every 5 yrs

Loads of applications!!!

sition (DMD), have been developed and extended in real-world applications. In this review, we provide a comprehensive overview of the theory, describing recent theoretical and computational advances and highlighting these methods with a diverse range of applications. We discuss the challenges in the rapidly growing field of machine learning and significantly transform the theory

Key words. dynamical systems, Koopman operator, data-driven theory, operator theory, dynamic mode decomposition

AMS subject classifications. 34A34, 37A30, 37C10, 37M10, 37M20, 37M25, 37M26, 37M27, 37M28, 37M29, 37M30, 37M31, 37M32, 37M33, 37M34, 37M35, 37M36, 37M37, 37M38, 37M39, 37M40, 37M41, 37M42, 37M43, 37M44, 37M45, 37M46, 37M47, 37M48, 37M49, 37M50, 37M51, 37M52, 37M53, 37M54, 37M55, 37M56, 37M57, 37M58, 37M59, 37M60, 37M61, 37M62, 37M63, 37M64, 37M65, 37M66, 37M67, 37M68, 37M69, 37M70, 37M71, 37M72, 37M73, 37M74, 37M75, 37M76, 37M77, 37M78, 37M79, 37M80, 37M81, 37M82, 37M83, 37M84, 37M85, 37M86, 37M87, 37M88, 37M89, 37M90, 37M91, 37M92, 37M93, 37M94, 37M95, 37M96, 37M97, 37M98, 37M99, 37M00

Cambridge, United Kingdom
 e-mail address: m.colbrook@damp.cam.ac.uk

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Outline

Question: *When can we reliably learn Koopman spectral properties from system data, and when is it impossible?*

- Constructing adversaries – *impossibility theorem* **(a general strategy)**
- Towers of algorithms – *possibility theorem* **(problem-specific algorithm)**
- The Solvability Complexity Index Hierarchy – *classifications* **(general tool)**
- Where does this leave us?

Theorem A (impossibility)

Implies \mathcal{K} is unitary



Class of systems: $\Omega_{\mathbb{D}} = \{F: \bar{\mathbb{D}} \rightarrow \bar{\mathbb{D}} \mid F \text{ cts, measure preserving, invertible}\}.$

Data an algorithm can use: $\mathcal{T}_F = \{(x, y_m) \mid x \in \bar{\mathbb{D}}, \|F(x) - y_m\| \leq 2^{-m}\}.$

Theorem A: There **does not exist** any sequence of deterministic algorithms $\{\Gamma_n\}$ using \mathcal{T}_F such that $\lim_{n \rightarrow \infty} \Gamma_n(F) = \text{Sp}(\mathcal{K}_F) \quad \forall F \in \Omega_{\mathbb{D}}.$

NB: Similarly, no random algorithms converging with probability $> 1/2$.

Proof idea: Constructing an adversary

$$F_0: \text{rotation by } \pi, \text{Sp}(\mathcal{K}_{F_0}) = \{\pm 1\}$$

Phase transition lemma: Let $X = \{x_1, \dots, x_N\}, Y = \{y_1, \dots, y_N\}$ be distinct points in annulus $\mathcal{A} = \{x \in \mathbb{D} \mid 0 < R < \|x\| < r < 1\}$ with $X \cap Y = \emptyset$. There exists a measure-preserving homeomorphism H such that H acts as the identity on $\mathbb{D} \setminus \mathcal{A}$ and $H(y_j) = F_0(H(x_j)), j = 1, \dots, N$.

Conjugacy of data ($x_j \rightarrow y_j$) with F_0

Idea: Use lemma to trick any algorithm into oscillating between spectra.

Proof idea: Constructing an adversary

Suppose (for contradiction) $\{\Gamma_n\}$ uses \mathcal{J}_F , $\lim_{n \rightarrow \infty} \Gamma_n(F) = \text{Sp}(\mathcal{K}_F) \quad \forall F \in \Omega_{\mathbb{D}}$.

Build an **adversarial** F ...

$$\mathcal{J}_F = \{(x, y_m) \mid \|F(x) - y_m\| \leq 2^{-m}\}$$

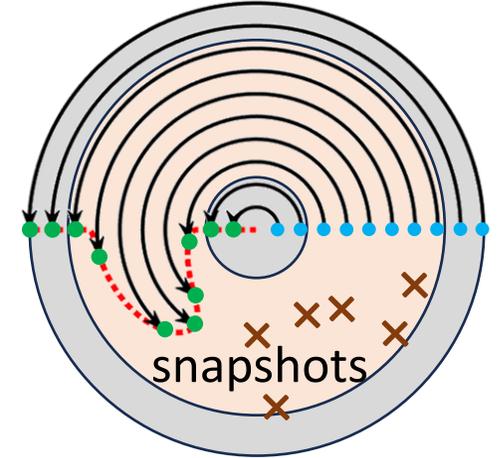
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Suppose (for contradiction) $\{\Gamma_n\}$ uses \mathcal{T}_F , $\lim_{n \rightarrow \infty} \Gamma_n(F) = \text{Sp}(\mathcal{K}_F) \quad \forall F \in \Omega_{\mathbb{D}}$.

Build an **adversarial** F ...

$$\widetilde{F}_1(r, \theta) = (r, \theta + \pi + \phi(r)), \text{supp}(\phi) \subset [1/4, 3/4]$$

$$\text{Sp}(\mathcal{K}_{\widetilde{F}_1}) = \mathbb{T} \text{ (unit circle).}$$



$$\mathcal{T}_F = \{(x, y_m) \mid \|F(x) - y_m\| \leq 2^{-m}\}$$

Proof idea: Constructing an adversary

Suppose (for contradiction) $\{\Gamma_n\}$ uses \mathcal{T}_F , $\lim_{n \rightarrow \infty} \Gamma_n(F) = \text{Sp}(\mathcal{K}_F) \quad \forall F \in \Omega_{\mathbb{D}}$.

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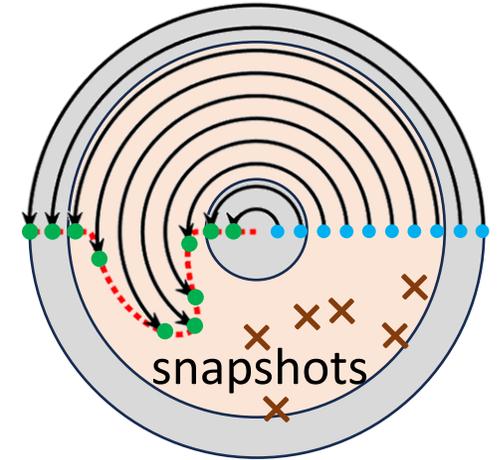
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$$\text{Sp}(\mathcal{K}_{\widetilde{F}_1}) = \mathbb{T} \text{ (unit circle).}$$

$$\lim_{n \rightarrow \infty} \Gamma_n(\widetilde{F}_1) = \text{Sp}(\mathcal{K}_{\widetilde{F}_1}) \Rightarrow \exists n_1 \text{ s.t. } \text{dist}(i, \Gamma_{n_1}(\widetilde{F}_1)) \leq 1.$$

BUT Γ_{n_1} uses finite amount of info to output $\Gamma_{n_1}(\widetilde{F}_1)$.

Let X, Y correspond to these snapshots.



$$\mathcal{T}_F = \{(x, y_m) \mid \|F(x) - y_m\| \leq 2^{-m}\}$$

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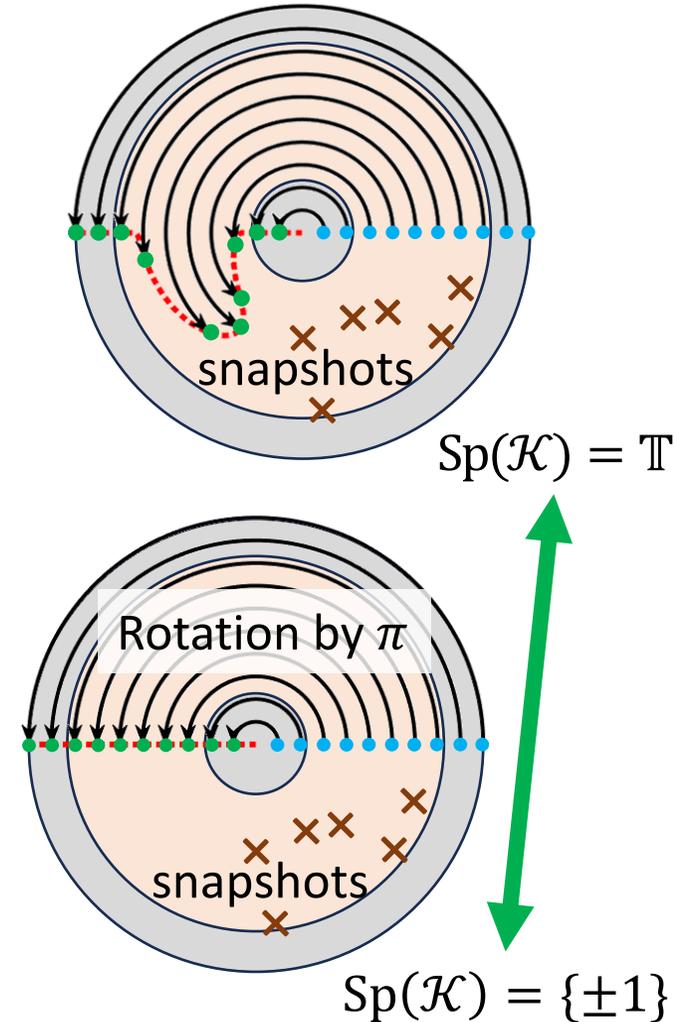
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Let X, Y correspond to these snapshots.

Lemma: $F_1 = H_1^{-1} \circ F_0 \circ H_1$ on annulus \mathcal{A}_1 .

Consistent data $\Rightarrow \Gamma_{n_1}(F_1) = \Gamma_{n_1}(\widetilde{F}_1)$, $\text{dist}(i, \Gamma_{n_1}(F_1)) \leq 1$

BUT $\text{Sp}(\mathcal{K}_{F_1}) = \text{Sp}(\mathcal{K}_{F_0}) = \{\pm 1\}$



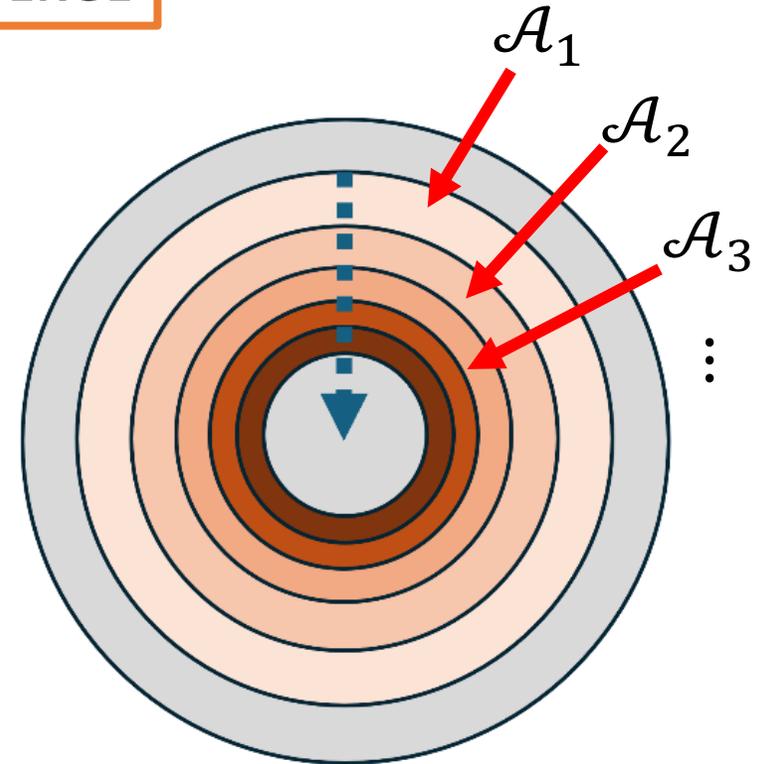
Proof idea: Constructing an adversary

Inductive step: Repeat on annuli, $F_k = H_k^{-1} \circ F_0 \circ H_k$ on \mathcal{A}_k . $F = \lim_{k \rightarrow \infty} F_k$

Consistent data $\Rightarrow \Gamma_{n_k}(F) = \Gamma_{n_k}(\widetilde{F}_k)$, $\text{dist}(i, \Gamma_{n_k}(F)) \leq 1$, $n_k \rightarrow \infty$

BUT $\text{Sp}(\mathcal{K}_F) = \text{Sp}(\mathcal{K}_{F_0}) = \{\pm 1\}$

CANNOT CONVERGE



Cascade of disks

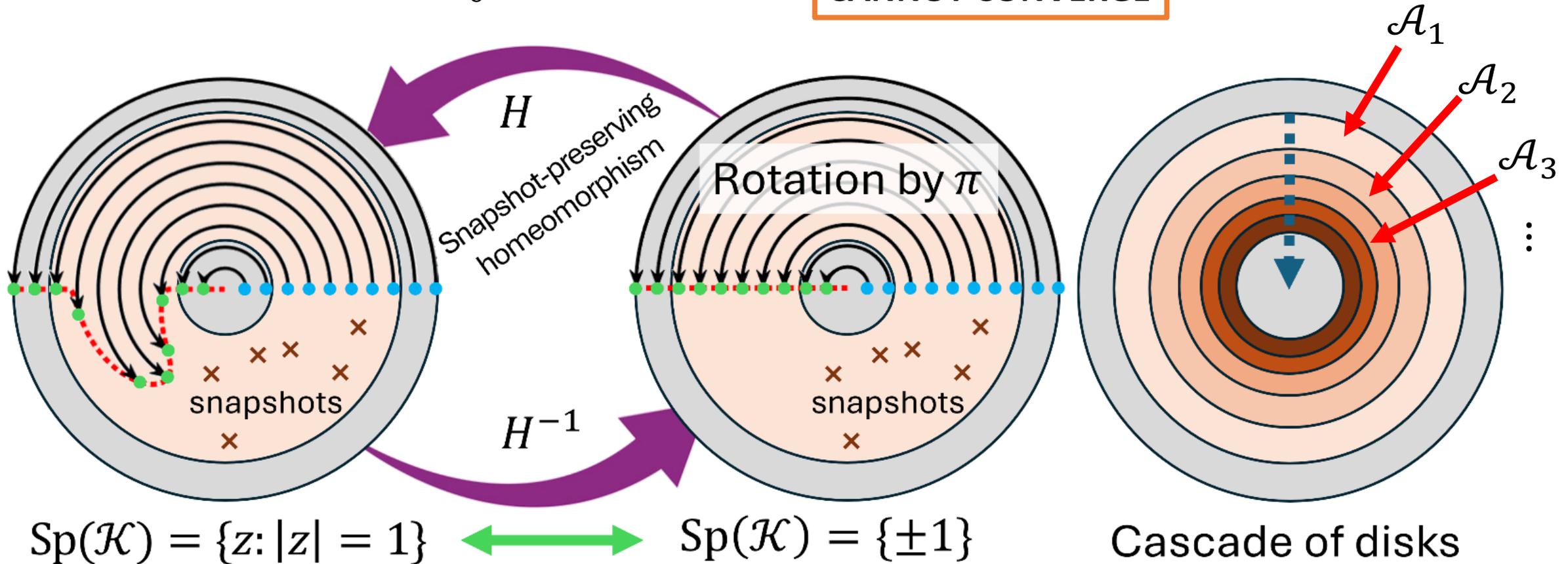
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BUT $\text{Sp}(\mathcal{K}_F) = \text{Sp}(\mathcal{K}_{F_0}) = \{\pm 1\}$

CANNOT CONVERGE



Theorem B (possibility)

$\Omega_{\mathcal{X}}^m = \{F: \mathcal{X} \rightarrow \mathcal{X} \mid F \text{ cts, measure preserving}\}.$

$\mathcal{T}_F = \{(x, y_m) \mid x \in \mathcal{X}, \|F(x) - y_m\| \leq 2^{-m}\}.$

Theorem B: There **exists deterministic** algorithms $\{\Gamma_{n_2, n_1}\}$ using input data \mathcal{T}_F such that $\lim_{n_2 \rightarrow \infty} \lim_{n_1 \rightarrow \infty} \Gamma_{n_2, n_1}(F) = \text{Sp}(\mathcal{K}_F) \quad \forall F \in \Omega_{\mathcal{X}}^m.$

Double limit $\lim_{n_2 \rightarrow \infty} \lim_{n_1 \rightarrow \infty}$

Proof sketch

- $\lim_{N \rightarrow \infty} \lim_{M \rightarrow \infty} \gamma_{N,M}(F, z) = \|(\mathcal{K}_F - zI)^{-1}\|^{-1}$.

N = size of basis, M = amount of data (quadrature).

Proof sketch

- $\lim_{N \rightarrow \infty} \lim_{M \rightarrow \infty} \gamma_{N,M}(F, z) = \|(\mathcal{K}_F - zI)^{-1}\|^{-1}$.

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- Measure-preserving $\Rightarrow \|(\mathcal{K}_F - zI)^{-1}\|^{-1} = \text{dist}(z, \text{Sp}(\mathcal{K}_F))$.

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- Measure-preserving $\Rightarrow \|(\mathcal{K}_F - zI)^{-1}\|^{-1} = \text{dist}(z, \text{Sp}(\mathcal{K}_F))$.
- Local N -adaptive minimisation of $\gamma_{N,M}(F, z)$ to approximate $\text{Sp}(\mathcal{K}_F)$

Proof sketch

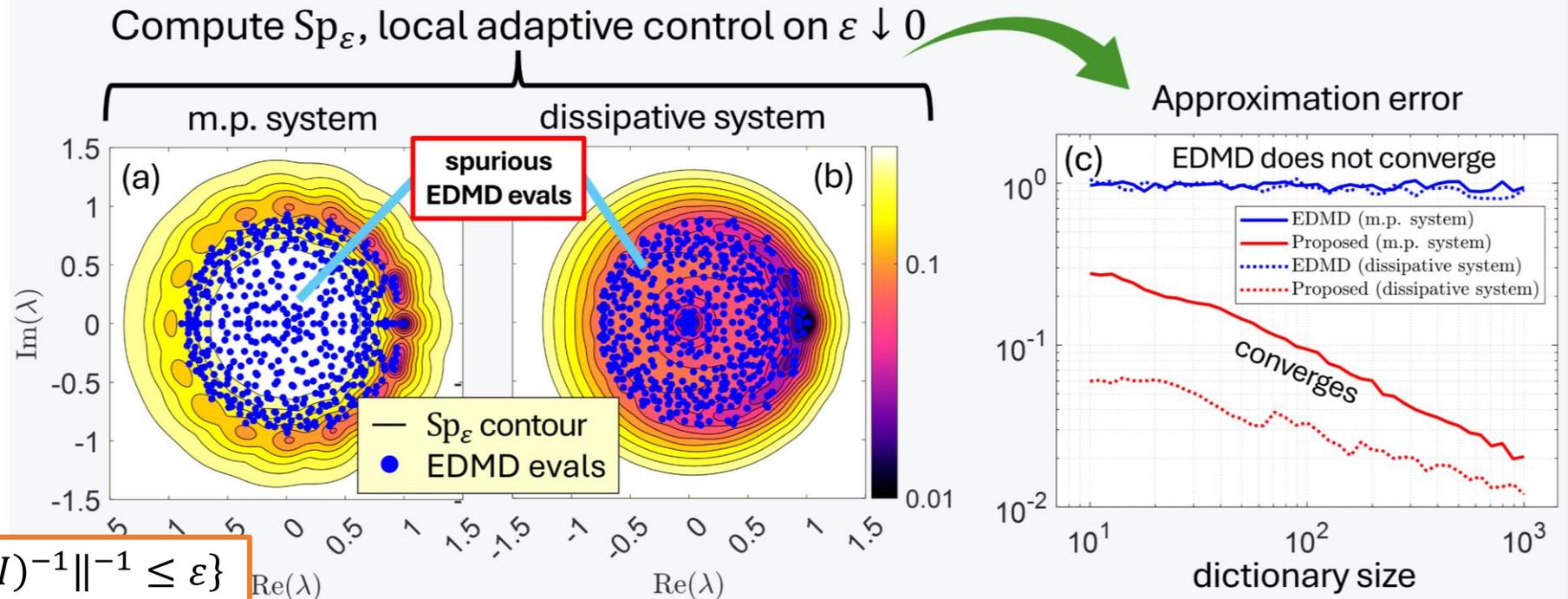
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- Local N -adaptive minimisation of $\gamma_{N,M}(F, z)$ to approximate $\text{Sp}(\mathcal{K}_F)$

“Extended Dynamic Mode Decomposition” is the gold-standard Galerkin method.



Limits of limits: Towers of algorithms

Def: $\{\Gamma_{n_k, \dots, n_1}\}$ with $\lim_{n_k \rightarrow \infty} \dots \lim_{n_1 \rightarrow \infty} \Gamma_{n_k, \dots, n_1}$ convergent a ***tower of algorithms***.

First appeared in dynamical systems theory: ← algorithms



Steve Smale

“Is there any purely iterative convergent rational map for polynomial zero finding?”



Curtis McMullen

“Yes for cubic, no for higher degree. Quartic and quintic can be solved using towers of algorithms. Sextic cannot be solved in any number of limits.”

- Smale, “On the efficiency of algorithms of analysis.” **Bull. Am. Math. Soc.**, 1985.
- McMullen, “Families of rational maps and iterative root-finding algorithms.” **Annals Math.**, 1987.
- McMullen, “Braiding of the attractor and the failure of iterative algorithms.” **Invent. Math.** 1988.
- Doyle, McMullen, “Solving the quintic by iteration.” **Acta Math.**, 1989.

Classifications: *Solvability Complexity Index (SCI)*

SCI: Fewest number of limits needed to solve a computational problem.

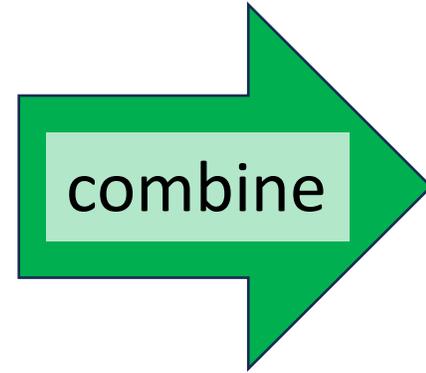
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Theorem A: $SCI > 1$

Theorem B: $SCI \leq 2$



$SCI = 2$

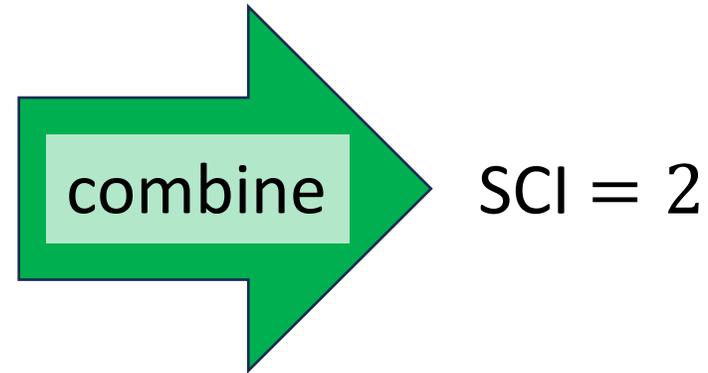
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So far literature has only proven upper bounds, that need not be sharp...

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Results from Koopman literature

SCI: Fewest number of limits needed to solve a computational problem.

Algorithm	Comments/Assumptions	Spectral Problem's Corresponding SCI Upper Bound			
		<i>KMD</i>	<i>Spectrum</i>	<i>Spectral Measure (if m.p.)</i>	<i>Spectral Type (if m.p.)</i>
Extended DMD [47]	general L^2 spaces	$\text{SCI} \leq 2^*$	N/C	N/C	n/a
Residual DMD [44]	general L^2 spaces	$\text{SCI} \leq 2^*$	$\text{SCI} \leq 3^*$	$\text{SCI} \leq 2^*$	varies, see [84] e.g., a.c. density: $\text{SCI} \leq 2^*$
Measure-preserving EDMD [45]	m.p. systems	$\text{SCI} \leq 1$	N/C	$\text{SCI} \leq 2^*$ (general) $\text{SCI} \leq 1$ (delay-embedding)	n/a
Hankel DMD [85]	m.p. ergodic systems	$\text{SCI} \leq 2^*$	N/C	N/C	n/a
Periodic approximations [86]	m.p. + ω a.c.	$\text{SCI} \leq 2$	N/C	$\text{SCI} \leq 2$ (see [87])	a.c. density: $\text{SCI} \leq 3$
Christoffel–Darboux kernel [40]	m.p. ergodic systems	$\text{SCI} \leq 3$	n/a	$\text{SCI} \leq 2$	e.g., a.c. density: $\text{SCI} \leq 2$
Generator EDMD [88]	cts.-time, samples ∇F (otherwise additional limit)	$\text{SCI} \leq 2$	N/C	$\text{SCI} \leq 2$ (see [89])	n/a
Compactification [42]	cts.-time, m.p. ergodic systems	$\text{SCI} \leq 4$	N/C	$\text{SCI} \leq 4$	n/a
Resolvent compactification [43]	cts.-time, m.p. ergodic systems	$\text{SCI} \leq 5$	N/C	$\text{SCI} \leq 5$	n/a
Diffusion maps [90] (see also [10])	cts.-time, m.p. ergodic systems	$\text{SCI} \leq 3$	n/a	n/a	n/a

Are these sharp?

Previous techniques prove upper bounds on SCI.

“N/C”: method need not converge. “n/a”: algorithm not applicable to problem.

Also in Ulam’s method for Markov processes, SRB measure computation, control,...

Classifications: *Solvability Complexity Index (SCI)*

SCI: Fewest number of limits needed to solve a computational problem.

- Δ_1 : One limit, full error control. E.g., $d(\Gamma_n(F), \text{Sp}(\mathcal{K}_F)) \leq 2^{-n}$.

- Δ_{m+1} : $\text{SCI} \leq m$.

- Σ_m : $\text{SCI} \leq m$, final limit from below.

E.g., Σ_1 : $\Gamma_n(F) \subset \text{Sp}(\mathcal{K}_F) + B_{2^{-n}}(0)$.

- Π_m : $\text{SCI} \leq m$, final limit from above.

E.g., Π_1 : $\text{Sp}(\mathcal{K}_F) \subset \Gamma_n(F) + B_{2^{-n}}(0)$.

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Classifications: *Solvability Complexity Index (SCI)*

SCI: Fewest number of limits needed to solve a computational problem.

- Δ_1 : One limit, full error control. E.g., $d(\Gamma_n(F), \text{Sp}(\mathcal{K}_F)) \leq 2^{-n}$.

- Δ_{m+1} : $\text{SCI} \leq m$.

- Σ_m : $\text{SCI} \leq m$, final limit from below.

E.g., Σ_1 : $\Gamma_n(F) \subset \text{Sp}(\mathcal{K}_F) + B_{2^{-n}}(0)$.

- Π_m : $\text{SCI} \leq m$, final limit from above.

E.g., Π_1 : $\text{Sp}(\mathcal{K}_F) \subset \Gamma_n(F) + B_{2^{-n}}(0)$.

verification

trust output

covers spectrum

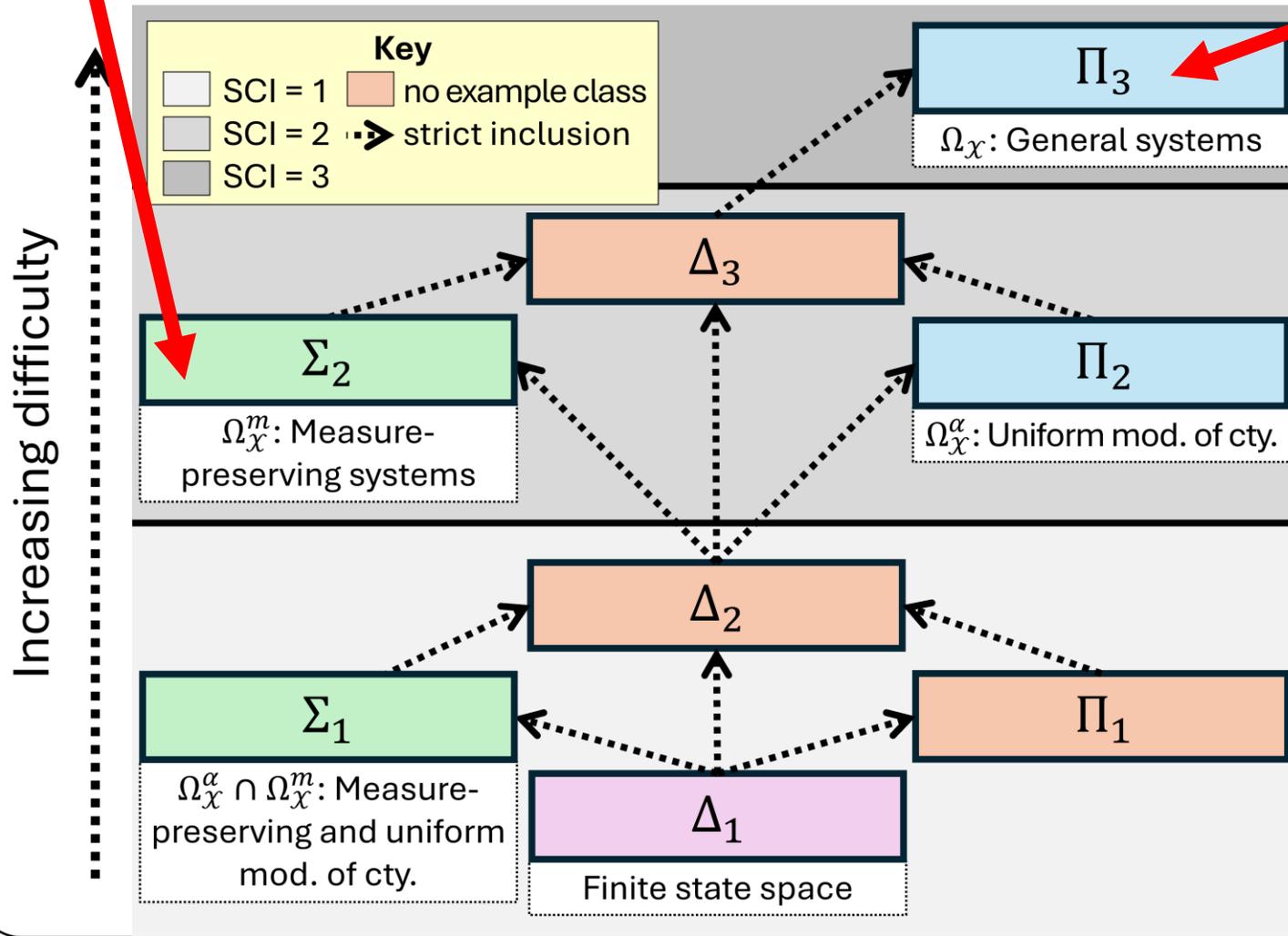
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Theorems A + B

Classification for Koopman I

3 limits needed in general!

SCI hierarchy of computing the spectrum



Different classes:

$$\Omega_X = \{F: X \rightarrow X \mid F \text{ cts}\}$$

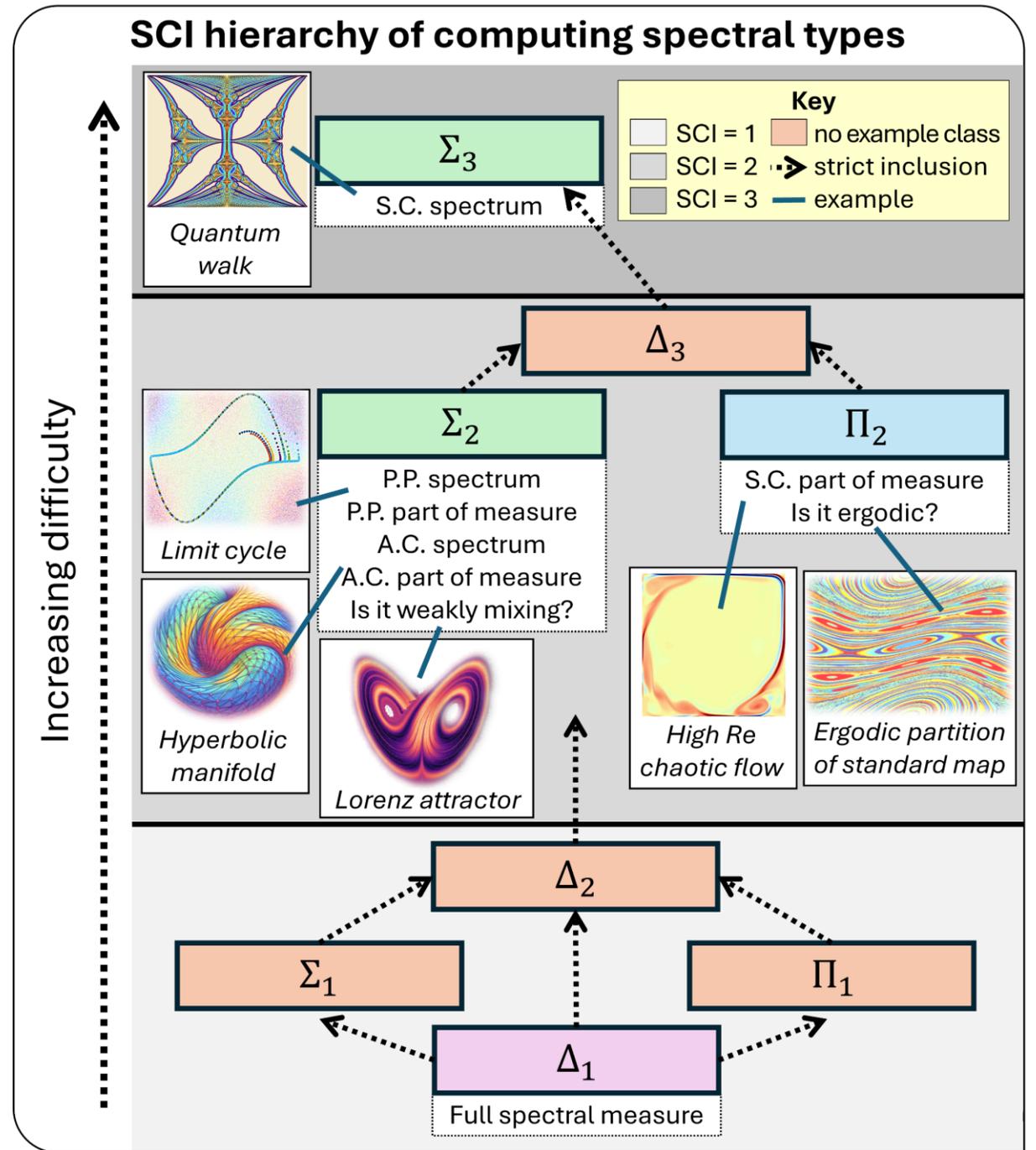
$$\Omega_X^m = \{F: X \rightarrow X \mid F \text{ cts, m. p.}\}$$

$$\Omega_X^\alpha = \{F: X \rightarrow X \mid F \text{ mod. cty. } \alpha\}$$

$$[d_X(F(x), F(y)) \leq \alpha(d_X(x, y))]$$

Optimal algorithms and classifications of dynamical systems.

Classification for Koopman II

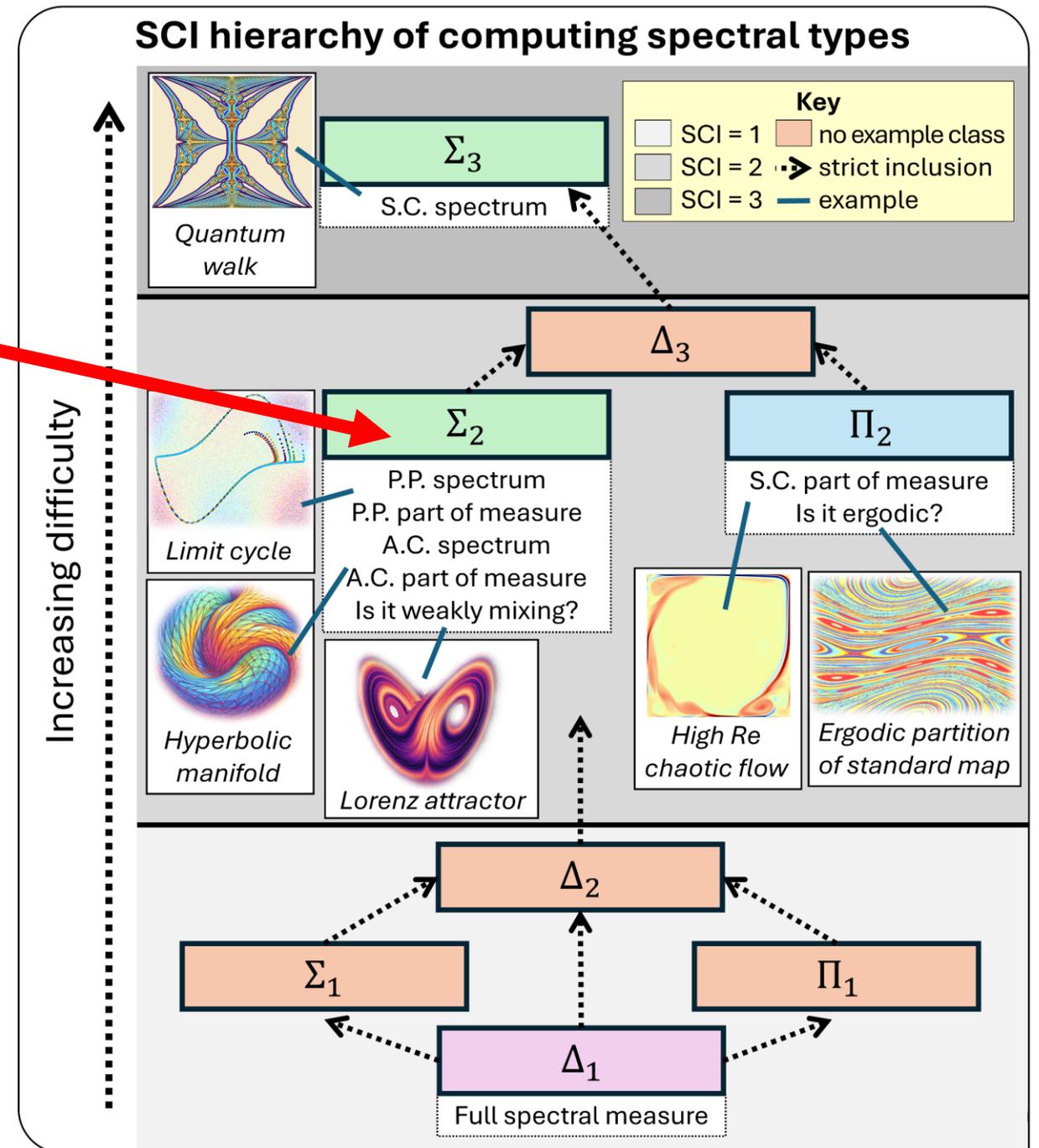


Classification for Koopman II

Example: Theorem C

For smooth, measure-preserving systems on a torus, learning eigenfunctions or even determining if there are any has $SCI = 2$ (even if we can sample derivatives).

Finding finite-dimensional embeddings in which the dynamics are linear (e.g., autoencoders, latent space representation) is very hard!



General tool in data-driven dynamical systems/PDEs

Adversarial arguments generalize to:

- “Learning the F ”. E.g., SINDy $(x_{n+1} = F(x_n))$
- Solving PDEs with neural networks (PINNs)
- Learning PDEs from forcing-solution pairs (e.g., hyperbolic)

-
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Where does this leave us?

- Many problems **NECESSARILY** require multiple limits.
- New tools for **lower bounds** (impossibility results) for Koopman learning.
- Combine with **upper bounds** (algorithms)
⇒ **classify difficulty** of problems + **prove optimality** of algorithms.
- Ergodic theory + approximation theory + computational analysis
⇒ started to map out this terrain.

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 - Other function spaces.
 - Partial observations, continuous-time.
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Where does your problem/method fit into the SCI hierarchy? Is it optimal?

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